

POWERING THE NEW ENGINEER TO TRANSFORM THE FUTURE

Department of Electrical & Computer Engineering

MeXT: A Flow for Multiprocessor Exploration

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Agenda

- Motivation
- The MeXT Design Flow
- MPI for Embedded Multiprocessor Systems
- Case Study
- Summary
- Future Work: Security and Reliability Extension



Motivation

- Increasing high computation demand
- Increasing Complexity due to multiple heterogeneous components in a system
- Heterogeneous systems are not optimized always for a set of applications
- An optimal communication among processing components is missing
- Parallelization required for some applications for performance increase



Motivation

Approach:

- Define a quality/tolerance factor for each component/task in the system
- Reformulate original optimization problem
- Provide an architecture that will enforce a high-level of reliability with minimum redundancy
- Provide methods to optimize architectural resources at run-time

Motivation

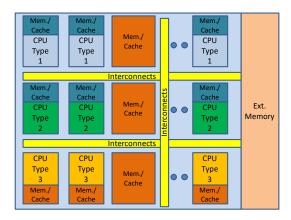
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Approach – Problem Reformulation

- Many problems are defined as on-line optimization problem, with the goal of producing a optimal output under environmental constraints
- Reformulate the problem as the production of the closest output under the same environmental constraints.
- Use current solver to devise an online solution
- However, provide a combined time-space redundancy that will improve the quality of results and keep the output close to the optimal result
- Use of reconfiguration logic, with extreme flexibility an performance to organize this work.
- Embedded Optimization to enforce rules at run-time

Target System

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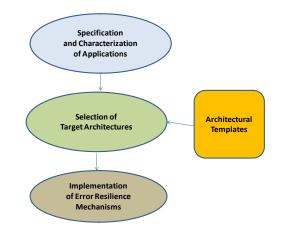
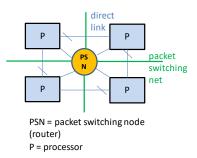
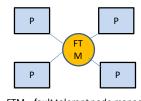
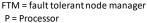
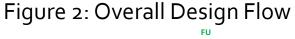


Figure 1: Target system architecture









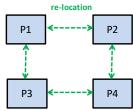


Figure 3: Communication Topologies

Target System

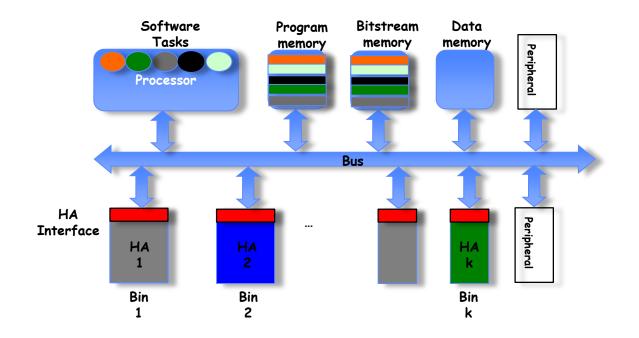


Figure 4: SoC Architecture with Hardware Accelerators

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MPSoC Exploration Design Flow

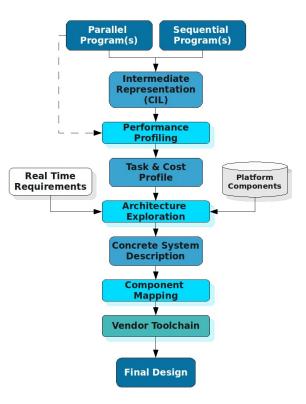
Input: Parallel programs and sequential programs, realtime constraints, platform components

- Intermediate representation of programs
 for performance profiling
- Cost optimization and architecture exploration
- Transforming of an Abstract Specification into a Platform-dependent Concrete Specification (e.g.: Platform 2 Xilinx ML310, CPU 2 PowerPC, Memory 2 BRAM, CommMedium 2 PLB,...)
- Generation of the platform-dependent hardware description files (Concrete Component Description)

Output: FPGA configuration file



• The *Concrete Specification* can also be the result of the architectural synthesis



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Semi-Automatic System Generation

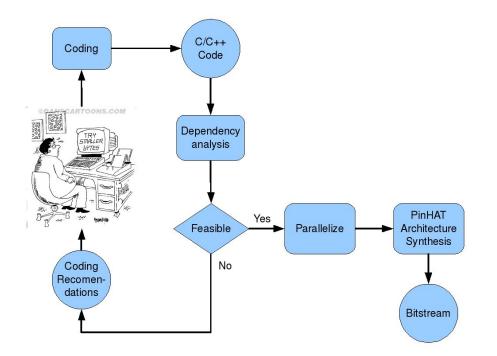


Figure 6: A Semi-Automatic System Generation Design Flow

Architecture Synthesis - Scope

Modeling goal:

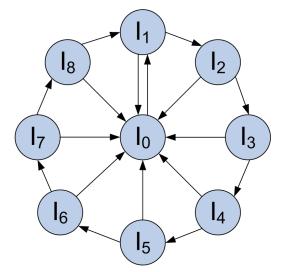
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- Allow selection of the best configuration (application-specific)
 - → Enable (runtime) optimization
- Lead to (automatic) generation of HW infrastructure
 - ➔ Cost models of HW components
 - → System model/description (PEs, networks, mapping, ...)
 - → SW configuration
 - ➔ Synthesis and device configuration
- Approach:
 - Synthesize an application-specific optimum system from parallel programs (high-level synthesis) using ILP



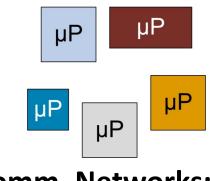
Architecture Synthesis - Problems

Parallel Program:

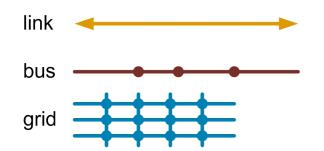


Problem: Simultaneous Task-Mapping and Selection of Resources

Processors:









MPSoC Design Objectives

- Construct an implementation with the desired functionality
- More challenging: simultaneously optimize numerous design metrics
- Metrics
 - Unit cost, NRE
 - Size, Power, a Weight
 - Performance
 - Flexibility, maintainability
 - Time-to-market, time-to-prototype

ILP Model

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Map all tasks

$$\sum_{j=0}^{m} x_{ij} = 1, \forall I_i$$

Max. usable FPGA area for PEs

$$v_j \le \sum_{i=0}^{m} x_{ij}, \forall J_j$$
$$A_{PE} \ge \sum_{j=0}^{m} v_{ij} \cdot a_j$$

 \boldsymbol{n}

Max # tasks per PE

$$\sum_{i=0}^{n} x_{ij} \cdot s_{ij} \le s_j, \forall I_i$$

Switching time for all tasks

$$T_{SWITCH} = \sum_{j=0}^{m} \sum_{i=0}^{n} x_{ij} \cdot t_j$$

ILP Model

Maximum Number of processes using a communication topology

$$\lambda_{i1,i2} = \frac{x_{i_1j_1} + x_{i_2j_2}}{2} \qquad \qquad y_k + \sum_{I_{i_1}, I_{i_2} \mid I_{i_1} \leqslant I_{i_2}} \lambda_{i1,i2} \le M_k, \forall C_k$$

Area Cost of Communication Network

$$A_{NET} \ge \sum_{j=0}^{m} A_k \cdot y_k$$

Total Area Cost

$$A \ge A_{PE} + A_{NET}$$

ILP Model

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Computation Time Cost for network topology

$$z_{ki_1i_2} \ge \lambda_{i1,i2} + y_k - 1$$

$$T_{NET} = \sum_{I_{i_1}, I_{i_2} | I_{i_1} \leqslant I_{i_2}} \left(\sum_{k=0}^K (D_{i_1 i_2} + \tau_k \cdot p_k) z_{k i_1 i_2} \right)$$

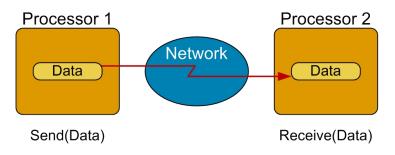
Total Computation Time Cost

$$T = \min\left(\sum_{j=0}^{m} \sum_{i=0}^{n} x_{ij} \cdot T_{ij} + T_{NET} + T_{SWITCH}\right)$$



Soc*MPI*

 Embedded Systems lack an efficient Message Passing Interface (MPI) communication library



- SocMPI on-Chip Communication library
- Compatible with a subset of the MPI Standard
- Hardware supported MPI functions (BCast, WTime)
- Configurable (MPI functions and networks)
- Memory requirements: 11 16 KByte

SocMPI

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- Two Layers: network dependent and network independent
- Currently supported : Direct Link, Star, Bus and Ring topology
- Easily extendable to other networks

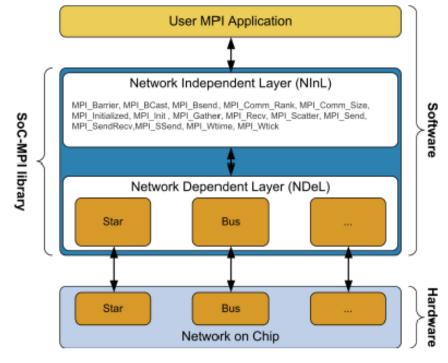


Figure 7: Layers of SoC-MPI

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Case Study 1: Mandelbrot Fractal

- Parallel implementation of Mandelbrot is used to analyze for generating profile.
- Three MPI tasks were allowed
- FSL direct-link, PLB bus and FSL-based bus used for exploration
- Simplified architecture is suggested in figure
- Table I shows parameter for available networks

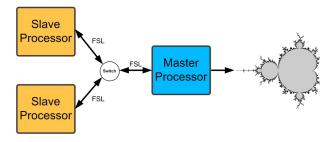


Figure 8: Simplified System

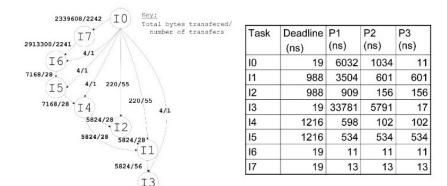
TABLE I PARAMETERS OF THE AVAILABLE NETWORKS

Name	Capacity	Area	Transfer Latency [μ s]	Arbitration Probability	Probability Bound	Туре
PLB	8	242	0.04	0.08	0.000016	bus
Star	8	213	0.04	0	0	star
FSL	1	21	0.02	0	0	link

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Case Study 2: WLAN 802.11g

- Parallel implementation of signal processing chain for IEEE 802.11g WLAN is used to analyze for generating profile.
- DAG shows communication pattern, task deadline and execution time
- Optimization is performed using MeXT and scheduling is analyzed
- Figure shows the resulting system architecture





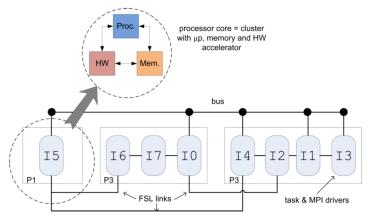


Figure 10: Architecture for WLAN program



Summary

- Design Systems Able to Cope with Probabilistic Behavior
- Framework for MPSoC Design Exploration
- Optimization for suitable Design
- Message Passing protocol for SoC-MPI



Future Work: Security and Resiliency Integration

- Secured Execution of Hardware/Software Threads in System on Chips
- Extending separation kernel policies within hardware components, and
- Enforcing corresponding access decision rules directly in hardware



Future Work: The Flask Security Architecture Model

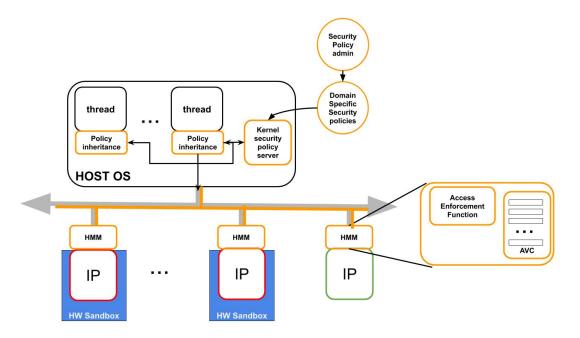
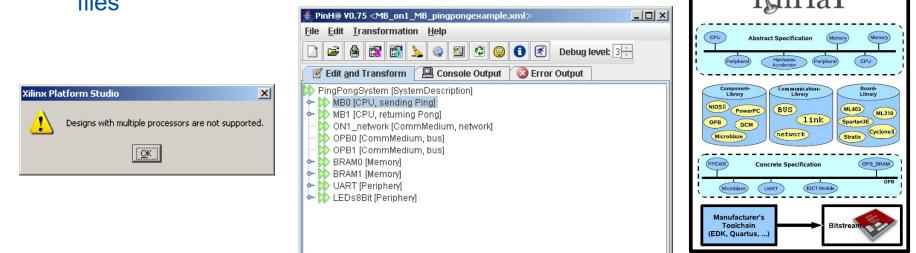


Figure 11: Flask Security Architecture with Hardware Management Module (HMM)



PinHaT (Plattform-independent <u>Ha</u>rdware Generation Tool)

- Funded By The German Research Foundation (DFG) 2005 -2010
 - Simplify the Design of MPSoCs
 - Vendor independent framework based on Java and XML
 - Abstract System Specification → Platform-dependent description files



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Security and Resiliency Integration

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Interface Integration

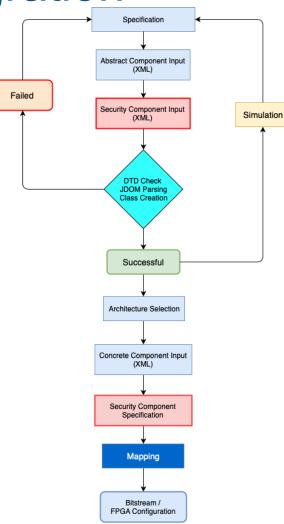
Security/Reliability Extension

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Security and Resiliency Integration

- The following extension has been added with PinHat
- Abstract security components (Hardware Sandbox and HMM) can be added now with existing abstract component input.
- While adding concrete component input, security components specification (e.g. size of HMM Access Vector Cache, I/O etc.) can be added.
- The communication interface between security components and bus can be specified











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